

(+2) mi baj van vele?

mind ortogonális a mintáihoz alkalmas

↑  
eddig egyszerű feltevések

$$y_c(k+1) = \text{sgn} \left\{ \sum_{j=1}^N w_{cj} y_j(k) \right\}$$

$l = \text{max}(N)$

$$w_{cj} = \frac{1}{N} \sum_{\alpha=1}^M s_i^\alpha s_j^\alpha$$

$$S^{-1} = \text{sgn} \{ \bar{W} S^{-1} \}$$

$$\bar{y}(0) = \bar{x}$$

$$M \leq \min \left\{ \frac{N}{2D}, \frac{N}{2 \log_2 N} \right\}$$

↑ capacity of the network

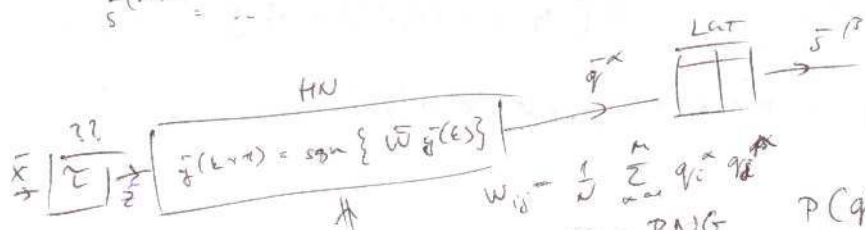
we have to choose these randomly

$$P(s_i^\alpha = 1) = P(s_i^\alpha = -1) = 0.5$$

↓  
RNG

$$\begin{aligned} \bar{s}^{(1)} &= (1 \ -1 \ -1 \ \dots \ 1) \\ \bar{s}^{(2)} &= (-1 \ 1 \ -1 \ -1 \ \dots \ 1) \\ &\vdots \\ \bar{s}^{(M)} &= \dots \end{aligned}$$

Stored memory patterns / prototypes are not random  
(they have some meaning)



$$\bar{q}^\alpha, \alpha = 1, \dots, M$$

↑  
H mapping

$$\bar{s}^\alpha, \alpha = 1, \dots, M$$

$$P(q_i^\alpha = 1) = P(q_i^\alpha = -1) = 0.5$$

not random but deterministically chosen which is dictated by the application

we map them

Topologically invariant mapping

$$\begin{aligned} d(\bar{x}, \bar{s}^\alpha) &< d(\bar{x}, \bar{s}^\beta) \quad \forall \alpha, \alpha \neq \beta, \alpha = 1, \dots, M \\ \Rightarrow d(\bar{z}, \bar{q}^\alpha) &< d(\bar{z}, \bar{q}^\beta) \quad \forall \alpha, \alpha \neq \beta, \alpha = 1, \dots, M \\ &\leq d(\bar{x}, \bar{s}^\alpha) \end{aligned}$$

Implementation of the topologically invariant mapping

$$z_l, z_c = \text{sgn} \left\{ \sum_{j=1}^N M_{lj} x_j \right\} \quad l = 1, \dots, L$$

$$M_{lj} := \frac{1}{N} \sum_{\alpha=1}^M q_l^\alpha \cdot s_j^\alpha$$

$$\bar{M} = \frac{1}{N} \sum_{\alpha=1}^M \bar{q}^\alpha \cdot \bar{s}^{\alpha T}$$

normalized inner prod.

$$\bar{z} = \text{sgn} \{ \bar{M} \cdot \bar{x} \}$$

↑  
matrix-vector mult.

$$a(\bar{x}, \bar{s}^\alpha) = \frac{1}{N} \sum_{i=1}^N x_i s_i^\alpha = \frac{1}{N} \bar{x}^T \bar{s}^\alpha \quad a(\bar{z}, \bar{q}^\alpha) = \frac{1}{N} \sum_{i=1}^N z_i q_i^\alpha$$

We have to prove that:  $a(\bar{x}, \bar{s}^n) > a(\bar{z}, \bar{q}^n)$

$$a(\bar{x}, \bar{s}^n) \leq a(\bar{z}, \bar{q}^n)$$

Es a skálázástól nagyobb,  
közelebb vanunk

Let's assume that:  $q_i^\beta = 1 \quad \forall i=1, \dots, N$

$$a(\bar{z}, \bar{q}^n) = \frac{1}{N} \sum_{i=1}^N z_i q_i^\beta = \frac{1}{N} \sum_{i=1}^N z_i \stackrel{n \rightarrow \infty}{\approx} E(z) = 1 \cdot P(z_i=1) + (-1) P(z_i=-1)$$

$$a(\bar{z}, \bar{q}^n) = P(z_i=1) - P(z_i=-1) =$$

$$= P(\text{sgn} \left\{ \sum_{j=1}^M M_{ij} x_j \right\} = 1) - P(\text{sgn} \left\{ \sum_{j=1}^M M_{ij} x_j \right\} = -1) =$$

$$= P(\text{sgn} \left\{ \sum_{j=1}^M \left( \sum_{\alpha=1}^M q_i^\alpha s_j^\alpha \right) x_j \right\} = 1) - P(\dots = -1) =$$

$$= P(\text{sgn} \left\{ \sum_{\alpha=1}^M q_i^\alpha \left( \frac{1}{N} \sum_{j=1}^N s_j^\alpha x_j \right) \right\} = 1) - P(\dots = -1) =$$

$$= P(\text{sgn} \left\{ q_i^\beta \frac{1}{N} \sum_{j=1}^N s_j^\beta x_j + \sum_{\substack{\alpha=1 \\ \alpha \neq \beta}}^M q_i^\alpha \frac{1}{N} \sum_{j=1}^N s_j^\alpha x_j \right\} = 1) - P(\dots = -1) =$$

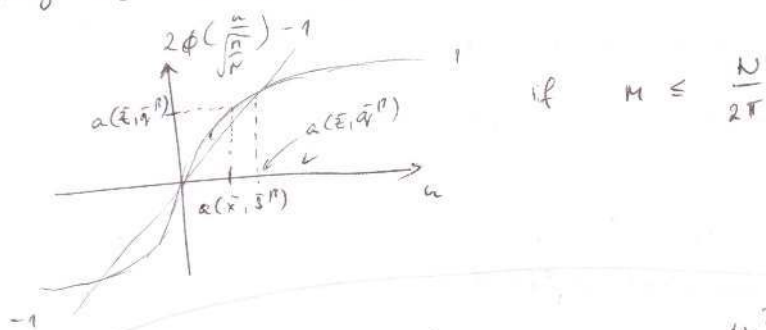
$$= P(\text{sgn} \left\{ a(\bar{x}, \bar{s}^n) + v_i \right\} = 1) - P(\text{sgn} \left\{ a(\bar{x}, \bar{s}^n) + v_i \right\} = -1) =$$

$$= P(a(\bar{x}, \bar{s}^n) + v_i \geq 0) - P(a(\bar{x}, \bar{s}^n) + v_i < 0)$$

$$= 2 \Phi \left( \frac{a(\bar{x}, \bar{s}^n)}{\sqrt{\frac{M}{N}}} \right) - 1$$

$$a(\bar{z}, \bar{q}^n) = 2 \Phi \left( \frac{a(\bar{x}, \bar{s}^n)}{\sqrt{\frac{M}{N}}} \right) - 1$$

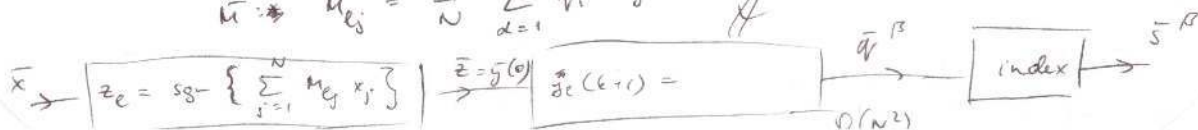
topologically invariant if  $a(\bar{z}, \bar{q}^n) \geq a(\bar{x}, \bar{s}^n)$



Given  $S = \{ s^\alpha, \alpha=1, \dots, M \}$   $M \leq \frac{N}{2 \log_2 N}$

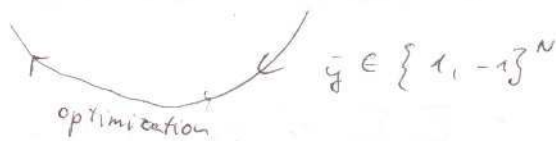
RNG  $Q = \{ \bar{q}^\alpha, \alpha=1, \dots, M \}$   $\bar{W} = \frac{1}{N} \sum_{\alpha=1}^M \bar{q}^\alpha \bar{q}^\alpha$

$\bar{u} = \{ u_{ij} = \frac{1}{N} \sum_{\alpha=1}^M q_i^\alpha s_j^\alpha \}$



Hopfield net as a combinatorial optimizer

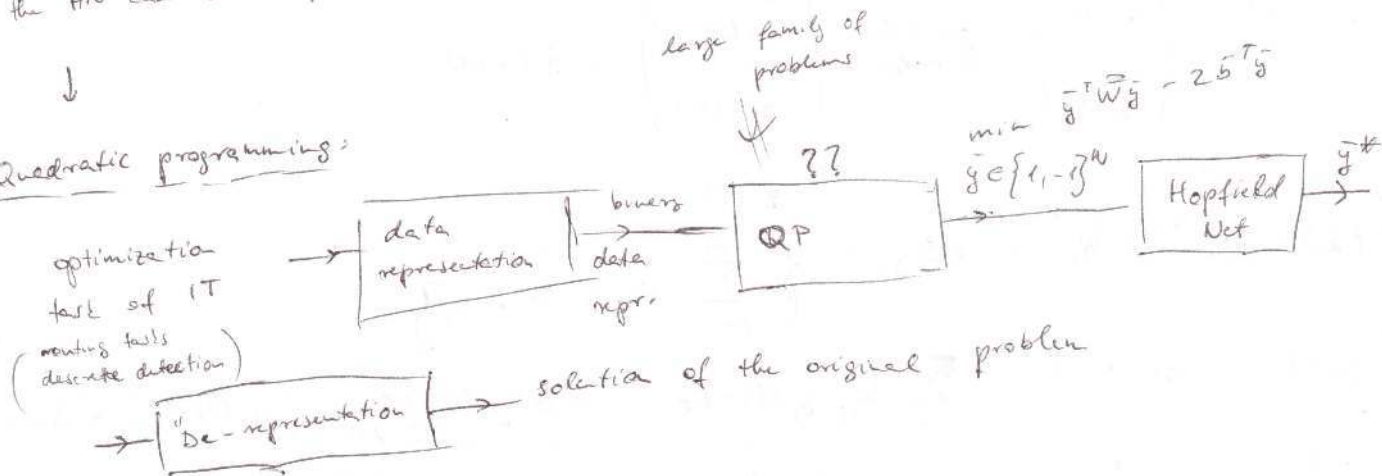
Intuitive description / idea:  $\bar{y}(k+1) = \text{sgn} \{ \bar{W} \bar{y}(k) - \bar{b} \}$   $\xrightarrow{\text{stability } O(N^2)}$   $\alpha(\bar{y}) = \bar{y}^T \bar{W} \bar{y} - 2 \bar{b}^T \bar{y}$



(if I have a test)

the HN can be interpreted as a discrete quadratic optimizer  $\leftarrow$  real-time

Quadratic programming:



~~we have to minimize~~ we have to minimize instead of maximizing

How to minimize a Quadratic form by HN?

Lemma:  $\alpha(\bar{y}) = \bar{y}^T \bar{W} \bar{y} - 2 \bar{b}^T \bar{y}$

$\bar{W} = W_{ij} = \begin{cases} w_{ij} & \text{if } i \neq j \\ 0 & \text{if } i = j \end{cases}$

$\alpha'(\bar{y}) = \bar{y}^T \bar{W}' \bar{y} - 2 \bar{b}^T \bar{y}$

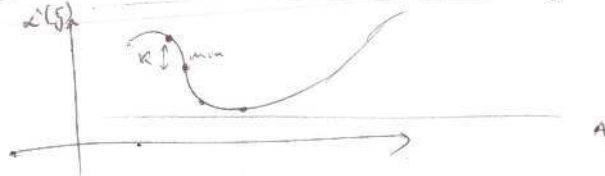
$\bar{y}^* = \min_{\bar{y} \in \{-1, 1\}^N} \alpha(\bar{y})$  ;  $\bar{y}^* = \min_{\bar{y} \in \{-1, 1\}^N} \alpha'(\bar{y})$

Proof:  $\alpha(\bar{y}) = \bar{y}^T \bar{W} \bar{y} - 2 \bar{b}^T \bar{y} = \sum_i \sum_j w_{ij} y_i y_j + 2 \sum_i b_i y_i =$

$= \sum_i w_{ii} y_i^2 + \sum_{i \neq j} \sum_j w_{ij} y_i y_j - 2 \sum_i b_i y_i = \sum_i w_{ii} + \sum_i \sum_j w'_{ij} y_i y_j - 2 \sum_i b_i y_i = \text{const} + \alpha'(\bar{y})$

$\Rightarrow$  their minimums are in the same place

$y_i(k+1) = -\text{sgn} \left\{ \sum_{j=1}^N W'_{ij} y_j(k) - b_i \right\} \rightarrow \min \alpha'(\bar{y})$



$TR \leq \frac{B-A}{K}$

$$B = N \|\bar{w}\| + 2\sqrt{N} \|\bar{b}\| \geq \alpha'(\bar{y})$$

$$A = -\bar{b}^T \bar{w}^T \bar{b} \leq \alpha'(\bar{y})$$

$$\Delta \alpha'(\bar{y}(k)) = \alpha'(\bar{y}(k+1)) - \alpha'(\bar{y}(k)) = \sum_c \sum_j w_{cj} y_j(k+1) y_c(k+1) - 2 \sum_c b_c y_c(k+1) =$$

$$- \sum_i \sum_j w_{ij}' y_i(k) y_j(k) + 2 \sum_i b_i y_i(k)$$

$$\bar{y}(k) = \begin{pmatrix} y_1(k) \\ y_2(k) \\ \vdots \\ y_n(k) \end{pmatrix} \xrightarrow{\ell = \text{mod } k} \begin{pmatrix} y_1(k) \\ y_2(k) \\ \vdots \\ y_c(k+1) \\ \vdots \\ y_n(k) \end{pmatrix} = \bar{y}(k+1)$$

$$\Delta \alpha'(\bar{y}) = \cancel{\Delta y_c^2(k) w_{cc}'} + 2 \Delta y_c(k) \left\{ \sum_{j=1}^N w_{cj} y_j(k) - b_c \right\}$$

$\frac{y_c(k)}{-1}$	$\frac{y_c(k+1)}{1}$	$\frac{\sum_j w_{cj}' y_j(k) - b_c}{\leq -k}$	$\frac{2 \Delta y_c(k) \left\{ \sum_{j=1}^N w_{cj}' y_j(k) - b_c \right\}}{-4k}$	$= \Delta \alpha(\ell)$
$1$	$-1$	$\geq k$	$-4k$	$2(-1 - (+1)) = -4$

↑  
always decreases

$$\text{If } \alpha(\bar{y}) = \bar{y}^T \bar{w} \bar{y} - 2\bar{b}^T \bar{y} \rightarrow \bar{w}' \rightarrow y_c(k+1) = -\text{sgn} \left\{ \sum_{j=1}^N w_{cj}' y_j(k) - b_c \right\}$$

Modified Hopfield Net can minimize quadratic forms over finite sets in P complexity.